The Descriptive Complexity of Bayesian Network Specifications

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- **1** A bit of motivation and context.
- 2 Setting up relational Bayesian network specifications.
- **3** The descriptive complexity results.
- 4 Conclusion: A model theory of Bayesian networks?

Bayesian networks, and repetitive patterns...



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Plates and Probabilistic Relational Models

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Logical Bayesian Networks (LBNs)

 $\begin{array}{ll} \mathsf{carrier}(x) \mid & \mathsf{founder}(x) \\ \mathsf{carrier}(x) \mid & \mathsf{mother}(m, x), \mathsf{father}(f, x), \mathsf{carrier}(m), \mathsf{carrier}(f) \\ \mathsf{suffers}(x) \mid & \mathsf{carrier}(x) \end{array}$

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Relational Bayesian Networks (RBNs)

Stan

model
$$\{y \sim std_normal();\}$$

 To study the connection between the language used to specify Bayesian networks and the expressivity of these Bayesian networks.

 To do so by resorting to descriptive complexity, a concept from finite model theory.











Non-root: $X \Leftrightarrow \phi$ where ϕ in a language \mathcal{L}

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$$\mathbb{P}(\operatorname{other}(\chi, y)) = 0.1.$$

$$fan(a) \quad fan(b) \quad fan(c)$$

$$\operatorname{es}(a, b) \quad \operatorname{likes}(a, c) \quad \operatorname{likes}(b, a) \quad \operatorname{likes}(b, c) \quad \operatorname{likes}(c, a) \quad \operatorname{likes}(c, b)$$

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A "fitness"-based model of asymmetric "likeship":

$$\begin{split} \mathbb{P}(\mathsf{fan}(\chi)) &= 0.2, \\ \mathbb{P}(\mathsf{likes}(\chi, y)) &\Leftrightarrow (\chi = y) \lor \\ & (\mathsf{fan}(\chi) \land \mathsf{fan}(y)) \lor \\ & \mathsf{other}(\chi, y), \\ \mathbb{P}(\mathsf{other}(\chi, y)) &= 0.1. \end{split}$$



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Since 1970, finite model theory, motivated by:

- applications: databases, complexity theory, language design.
- A central theme:

The connection between a language and what it can/cannot express.



Fagin's theorem (ESO captures NP):

A set of strings is in NP if and only if the strings are the finite models of a formula in *existential second-order logic*.

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$$\exists \mathsf{r},\mathsf{g},\mathsf{b}: \left(\begin{array}{c} (\forall \chi:(\text{ either } \mathsf{r}(\chi) \text{ or } \mathsf{g}(\chi) \text{ or } \mathsf{b}(\chi))) \land & \\ \\ \forall \chi, y: \mathsf{edge}(\chi, y) \to \left(\begin{array}{c} \neg(\mathsf{r}(\chi) \land \mathsf{r}(y)) \land \neg(\mathsf{g}(\chi) \land \mathsf{g}(y)) \\ \land \neg(\mathsf{b}(\chi) \land \mathsf{b}(y)) \end{array}\right) \right)$$

- A propositional Bayesian network encodes *a fixed* distribution.
- We can do more with a relational Bayesian network specification. How much more?

The complexity class PP

- Consider a *probabilistic* Turing machine:
 - probabilities are assigned to transitions.
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- The class of sets of strings thus decided within polynomial time is *exactly* PP.
- So: PP reflects the complexity of computing probabilities for a phenomenon that can be simulated by a polynomial probabilistic Turing machine.

What do relational Bayesian networks specifications capture?

Theorem: Bayesian networks in FFFO "capture" PP

A set of strings is in PP ${ m if}$ and only if the strings encode the domains/queries with probability > 1/2 for a relational Bayesian network specification.

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- That is: if you have a phenomenon that can be simulated by a probabilistic Turing machine, it can be modeled by a relational Bayesian network specification.
- More complex phenomena *cannot* be modeled with such a Bayesian network specification (unless complexity classes collapse!).

Suppose we allow second-order quantification in the specification. For instance:

$$\begin{array}{ll} \mathsf{partitioned} & \Leftrightarrow & \exists \mathsf{partition} : \forall \chi : \forall y : (\mathsf{edge}(\chi, y) \Rightarrow \\ & (\mathsf{partition}(\chi) \Leftrightarrow \neg \mathsf{partition}(y))) \, . \end{array}$$

Theorem: Bayesian networks in ESO "capture" PP^{NP}

A set of strings is in PP^{NP} if and only if

the strings encode the domains/queries with probability >1/2 for an existential second-order Bayesian network specification.

- A framework in which to study the expressivity of Bayesian network specifications.
 - A theory of descriptive complexity for Bayesian networks.
 - A "feasible" fragment of finite *probabilistic* model theory?

 In short: Relational specifications capture PP; existential second-order specifications capture PP^{NP}.