# The Descriptive Complexity of Bayesian Network Specifications 

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## Overview

1 A bit of motivation and context.
2 Setting up relational Bayesian network specifications.
3 The descriptive complexity results.
4 Conclusion: A model theory of Bayesian networks?

## Bayesian networks, and repetitive patterns...



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## Plates and Probabilistic Relational Models


(BUGS)

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## And also textual languages...

Logical Bayesian Networks (LBNs)
$\operatorname{carrier}(x) \mid$ founder $(x)$
carrier $(x) \mid$ mother $(m, x)$, father $(f, x)$, carrier $(m), \operatorname{carrier}(f)$
suffers $(x) \mid \operatorname{carrier}(x)$

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## Relational Bayesian Networks (RBNs)

$$
\begin{aligned}
\operatorname{burglary}(v) & =0.005 \\
\text { alarm }(v) & =(\operatorname{burglary}(v): 0.95,0.01) ; \\
\operatorname{calls}(v, w) & =(\text { neighbor }(v, w):(\operatorname{alarm}(w): 0.9,0.05), 0) ; \\
\text { alarmed }(v) & =\text { NoisyOr }\{\operatorname{calls}(w, v) \mid w: \text { neighbor }(w, v)\}
\end{aligned}
$$

## Stan

$$
\text { model } \quad\{y \sim \text { std_normal( }) ;\}
$$

## The goal here:

- To study the connection between the language used to specify Bayesian networks and the expressivity of these Bayesian networks.
- To do so by resorting to descriptive complexity, a concept from finite model theory.


## A framework, for propositional languages (Poole 1993)



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Non-root: $X \Leftrightarrow \phi$ where $\phi$ in a language $\mathcal{L}$

## Now, consider relational Bayesian network specifications

A "fitness"-based model of asymmetric "likeship":

$$
\begin{aligned}
\mathbb{P}(\operatorname{fan}(x))= & 0.2 \\
\mathbb{P}(\operatorname{likes}(x, y)) \Leftrightarrow & (x=y) \vee \\
& (\operatorname{fan}(x) \wedge \operatorname{fan}(y)) \vee \\
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$\operatorname{fan}(a) \quad \operatorname{fan}(b) \quad \operatorname{fan}(c)$
likes $(a, b)$ likes $(a, c)$ likes $(b, a)$ likes $(b, c)$ likes $(c, a)$ likes $(c, b)$
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## Model Theory:

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■ Since 1970, finite model theory, motivated by:

- applications: databases, complexity theory, language design.
- A central theme:

The connection between a language and what it can/cannot express.


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Input (string) $\longrightarrow$ Turing machine $\longrightarrow$ YES/NO

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Example: give relation edge, and test model for 3-colorability.
$\exists \mathrm{r}, \mathrm{g}, \mathrm{b}:\binom{(\forall x:($ either $\mathrm{r}(x)$ or $\mathrm{g}(x)$ or $\mathrm{b}(x))) \wedge}{\forall x, y: \operatorname{edge}(x, y) \rightarrow\left(\begin{array}{c}\neg(\mathrm{r}(x) \wedge \mathrm{r}(y)) \wedge \neg(\mathrm{g}(x) \wedge \mathrm{g}(y)) \\ \wedge \neg(\mathrm{b}(x) \wedge \mathrm{b}(y))\end{array}\right.}$

## So, let's look at Bayesian network specifications...

■ A propositional Bayesian network encodes a fixed distribution.

- We can do more with a relational Bayesian network specification. How much more?


## The complexity class PP

- Consider a probabilistic Turing machine:
- probabilities are assigned to transitions.
- Suppose that, for any input, we get to know (magically) whether the probability of accepting the input is larger than $1 / 2$.


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- Suppose that, for any input, we get to know (magically) whether the probability of accepting the input is larger than $1 / 2$.

■ The class of sets of strings thus decided within polynomial time is exactly PP.

- So: PP reflects the complexity of computing probabilities for a phenomenon that can be simulated by a polynomial probabilistic Turing machine.


## What do relational Bayesian networks specifications capture?

## Theorem: Bayesian networks in FFFO "capture" PP

A set of strings is in PP
if and only if
the strings encode the domains/queries with probability $>1 / 2$
for a relational Bayesian network specification.

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- That is: if you have a phenomenon that can be simulated by a probabilistic Turing machine, it can be modeled by a relational Bayesian network specification.
■ More complex phenomena cannot be modeled with such a Bayesian network specification (unless complexity classes collapse!).


## Going up to second order

- Suppose we allow second-order quantification in the specification. For instance:
partitioned $\Leftrightarrow \exists$ partition: $\forall x: \forall y:(\operatorname{edge}(x, y) \Rightarrow$ (partition $(x) \Leftrightarrow \neg$ partition $(y)))$.


## Theorem: Bayesian networks in ESO "capture" PPNP

A set of strings is in $\mathrm{PP}^{\mathrm{NP}}$ if and only if
the strings encode the domains/queries with probability $>1 / 2$ for an existential second-order Bayesian network specification.

■ A framework in which to study the expressivity of Bayesian network specifications.

- A theory of descriptive complexity for Bayesian networks.
- A "feasible" fragment of finite probabilistic model theory?

■ In short: Relational specifications capture PP; existential second-order specifications capture $\mathrm{PP}^{N P}$.

