

(Irrelevant) natural extension of choice functions

Rejection functions on gambles

Gambles The random variable X takes values in the finite possibility space \mathcal{X} . Any real-valued function on \mathcal{X} is called a **gamble**, and we collect all of them in $\mathcal{L}(\mathcal{X})$ (or \mathcal{L}). Given two gambles f and g in \mathcal{L} , we say that $f \leq g$ if $(\forall x \in \mathcal{X}) f(x) \leq g(x)$. Its strict variant $<$ on \mathcal{L} is given by: $f < g \Leftrightarrow (f \leq g \text{ and } f \neq g)$; we collect all f such that $0 < f$ in $\mathcal{L}_{>0}$.

We define $\mathcal{Q} \subseteq \mathcal{P}(\mathcal{L})$ as the collection of **non-empty but finite subsets** of \mathcal{L} .

Rejection function A **rejection function** R is a map

$$R: \mathcal{Q} \rightarrow \mathcal{Q} \cup \{\emptyset\}: A \mapsto R(A) \text{ such that } R(A) \subseteq A.$$

Rationality axioms We call a rejection function R on \mathcal{Q} **coherent** if for all A, A_1 and A_2 in \mathcal{Q} , all f and g in \mathcal{L} , and all λ in $\mathbb{R}_{>0}$:

- R₁. $R(A) \neq A$; [avoiding complete rejection]
- R₂. if $f < g$ then $f \in R(\{f, g\})$; [dominance]
- R₃. a. if $A_1 \subseteq R(A_2)$ and $A_2 \subseteq A$ then $A_1 \subseteq R(A)$; [Sen's α]
- b. if $A_1 \subseteq R(A_2)$ and $A \subseteq A_1$ then $A_1 \setminus A \subseteq R(A_2 \setminus A)$; [Aizerman]
- R₄. a. if $A_1 \subseteq R(A_2)$ then $\lambda A_1 \subseteq R(\lambda A_2)$; [scaling invariance]
- b. if $A_1 \subseteq R(A_2)$ then $A_1 + \{f\} \subseteq R(A_2 + \{f\})$. [independence]

We collect all coherent rejection functions in the set $\overline{\mathcal{R}}$.

The 'is not more informative than' relation Given two rejection functions R_1 and R_2 :

$$R_1 \text{ is not more informative than } R_2 \Leftrightarrow (\forall A \in \mathcal{Q})(R_1(A) \subseteq R_2(A)).$$

For any collection \mathbf{R} of rejection functions, its infimum is the rejection function given by

$$(\inf \mathbf{R})(A) := \bigcap \mathbf{R}(A) \text{ for all } A \text{ in } \mathcal{Q}.$$

If \mathbf{R} consists of coherent rejection functions, then $\inf \mathbf{R}$ is coherent itself.

Assessment Mostly, if a subject assesses his rejection functions, he will only provide an **incomplete specification**. He will state

"I assess $f \in R(B)$ for some B in \mathcal{Q} and f in B ."

or, if we assume that this assessment satisfies Axiom R₄b, equivalently:

"I assess $0 \in R(B)$ for some B in $\mathcal{Q}^0 := \{A \in \mathcal{Q} : 0 \in A\}$."

Formally, his assessment \mathcal{B} is a subset of \mathcal{Q}^0 :

Assessing $\mathcal{B} \subseteq \mathcal{Q}^0$ means: "my rejection function satisfies $(\forall B \in \mathcal{B}) 0 \in R(B)$ ".

Extending an assessment Given any assessment $\mathcal{B} \subseteq \mathcal{Q}^0$ and any rejection function R on \mathcal{Q} , we say that R **extends** the assessment \mathcal{B} if $B \in \mathcal{B} \Rightarrow 0 \in R(B)$ for every B in \mathcal{Q} .

Natural extension

Definition Given any assessment $\mathcal{B} \subseteq \mathcal{Q}^0$, the **natural extension** of \mathcal{B} is the rejection function

$$\mathcal{E}(\mathcal{B}) := \inf\{R \in \overline{\mathcal{R}} : (\forall B \in \mathcal{B}) 0 \in R(B)\},$$

where we let $\inf \emptyset$ be equal to $\text{id}_{\mathcal{Q}}$, the identity rejection function that maps every option set to itself.

A special rejection function The definition above is not so useful: it provides no explicit expression. To remedy this, consider the **special rejection function** $R_{\mathcal{B}}$ defined as:

$$R_{\mathcal{B}}(A) := \left\{ f \in A : (\exists A' \in \mathcal{Q})(A' \supseteq A \text{ and } (\forall g \in \{f\} \cup A' \setminus A) (A' - \{g\} \cap \mathcal{L}_{>0} \neq \emptyset \text{ or } (\exists B \in \mathcal{B}, \exists \mu \in \mathbb{R}_{>0}) \{g\} + \mu B \preceq A')) \right\}$$

for all A in \mathcal{Q} , where we define \preceq on \mathcal{Q} as:

$$A_1 \preceq A_2 \Leftrightarrow (\forall f_1 \in A_1)(\exists f_2 \in A_2) f_1 \leq f_2 \text{ for all } A_1 \text{ and } A_2 \text{ in } \mathcal{Q}.$$

Assessments avoiding complete rejection We say that $\mathcal{B} \subseteq \mathcal{Q}^0$ **avoids complete rejection** when $R_{\mathcal{B}}$ satisfies Axiom R₁.

Theorem 1. Consider any assessment $\mathcal{B} \subseteq \mathcal{Q}^0$. Then the following statements are equivalent:

- (i) \mathcal{B} avoids complete rejection;
- (ii) There is a coherent extension of \mathcal{B} : $(\exists R \in \overline{\mathcal{R}})(\forall B \in \mathcal{B}) 0 \in R(B)$;
- (iii) $\mathcal{E}(\mathcal{B}) \neq \text{id}_{\mathcal{Q}}$;
- (iv) $\mathcal{E}(\mathcal{B}) \in \overline{\mathcal{R}}$;
- (v) $\mathcal{E}(\mathcal{B})$ is the least informative rejection function that is coherent and extends \mathcal{B} .

When any (and hence all) of these equivalent statements hold, then $\mathcal{E}(\mathcal{B}) = R_{\mathcal{B}}$.

Application: purely binary assessments

Assume that the assessment $\mathcal{B} \subseteq \mathcal{Q}^0$ consist of only **binary sets**: $\mathcal{B} \subseteq \{\{0, f\} : f \in \mathcal{L}\}$. Therefore, $B := \bigcap \mathcal{B} \setminus \{0\} \subseteq \mathcal{L}$ is its corresponding **desirability assessment**.

Avoiding non-positivity Given any desirability assessment $B \subseteq \mathcal{L}$, we say that B **avoids non-positivity** when $\text{posi}(B) \cap \mathcal{L}_{\leq 0} = \emptyset$.

The inference mechanism for choice functions has the inference mechanism for desirability as a special case:

Theorem 2. Consider any purely binary assessment $\mathcal{B} \subseteq \mathcal{Q}^0$. Then $B := \bigcap \mathcal{B} \setminus \{0\} \subseteq \mathcal{L}$ avoids non-positivity if and only if \mathcal{B} avoids complete rejection, and if this is the case, then $\mathcal{E}(\mathcal{B}) = R_{\text{posi}(\mathcal{L}_{>0} \cup B)}$.

Proposition. Consider $\mathcal{B} \subseteq \mathcal{Q}^0$. If there is a coherent set of desirable gambles D such that $(\forall B \in \mathcal{B}) B \cap D \neq \emptyset$, then \mathcal{B} avoids complete rejection.

Therefore, this is a sufficient condition for avoiding complete rejection that is easy to check.

Binary choice

More-than-binary choice Rejection functions are **more-than-binary comparisons** of gambles. Given any rejection function R , we can summarise its binary behaviour in

$$D_R := \{f \in \mathcal{L} : 0 \in R(0, f)\};$$

if R is coherent, then D_R is a coherent set of desirable gambles.

Binary choice There might be multiple rejection functions associated to D ; the least informative one is

$$R_D(A) := \{f \in A : (\exists g \in A) g - f \in D\}$$

for all A in \mathcal{Q} . If D is coherent, then so is R_D . For any collection \mathcal{D}' of coherent sets of desirable gambles, we let $R_{\mathcal{D}'} := \inf\{R_D : D \in \mathcal{D}'\}$. Then

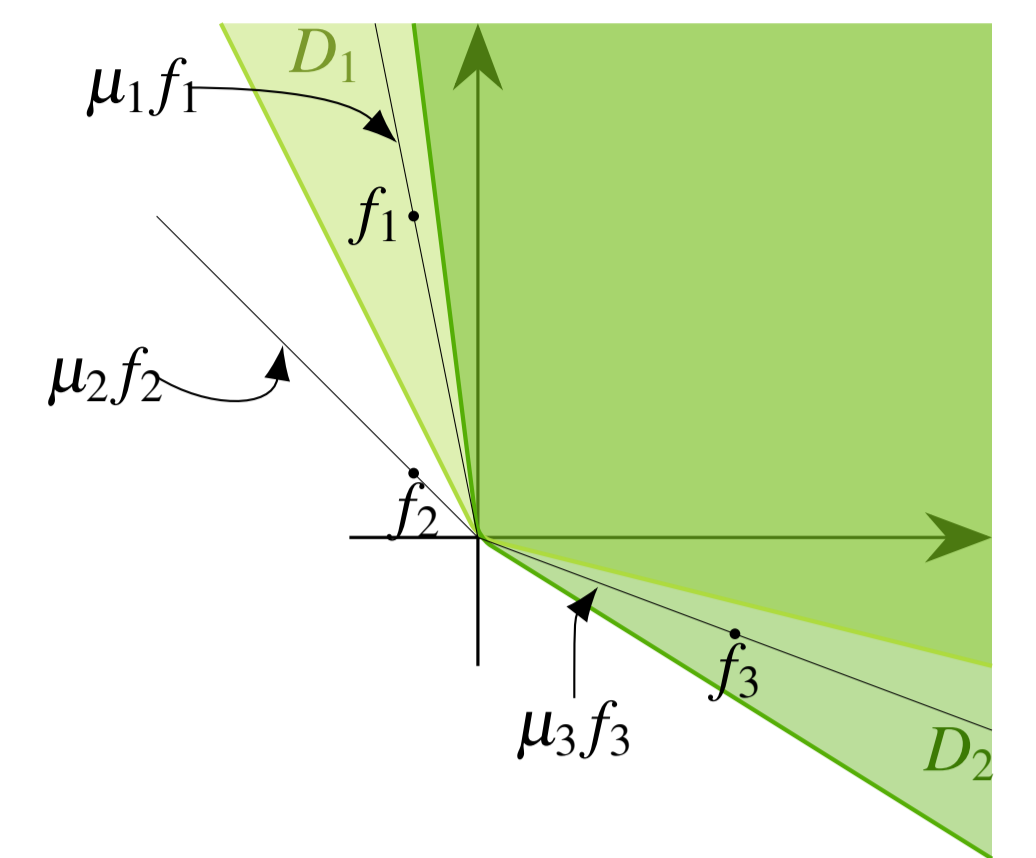
$$0 \in R_{\mathcal{D}'}(A \cup \{0\}) \Leftrightarrow (\forall D \in \mathcal{D}') D \cap A \neq \emptyset$$

for all A in \mathcal{Q} .

The natural extension of a desirability assessment $B \subseteq \mathcal{L}$ that avoids non-positivity, is

$$\mathcal{E}_{\mathcal{D}'}(B) := \text{posi}(\mathcal{L}_{>0} \cup B).$$

Example: infimum of binary choice

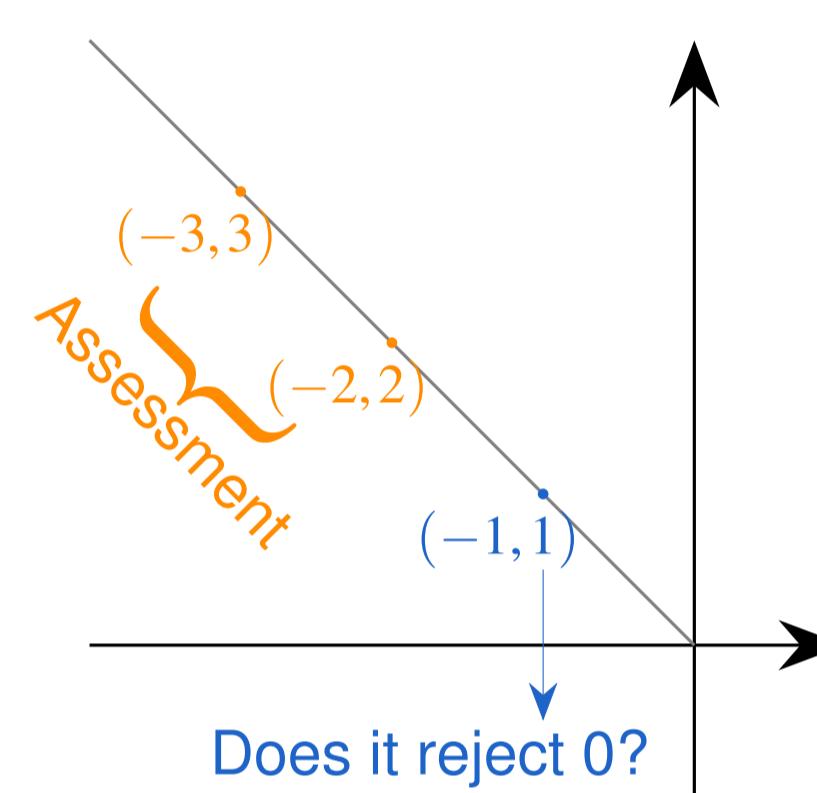


$\mathcal{D}' := \{D_1, D_2\}$ and $A := \{0, f_1, f_2, f_3\}$, so clearly $0 \in R_{\mathcal{D}'}(A)$, since $f_1 \in D_1$ and $f_3 \in D_2$.

Proposition. Consider any collection \mathcal{D}' of coherent sets of desirable gambles, any f_1, \dots, f_n in \mathcal{L} , and any $\mu_1 > 0, \dots, \mu_n > 0$. Then

$$0 \in R_{\mathcal{D}'}(\{0, f_1, \dots, f_n\}) \Leftrightarrow 0 \in R_{\mathcal{D}'}(\{0, \mu_1 f_1, \dots, \mu_n f_n\}).$$

Example: intrinsic non-binary choice



Assessment Consider the single assessment

$$\mathcal{B} := \{B\} \text{ where } B := \{0, (-2, 2), (-3, 3)\}.$$

It avoids complete rejection, by the Proposition in the frame **Application: purely binary assessments**. Therefore, $R_{\mathcal{B}}$ is a **coherent** rejection function.

Intrinsic non-binary choice

Note that $0 \in R_{\mathcal{B}}(0, (-2, 2), (-3, 3))$. We find that

$$0 \notin R_{\mathcal{B}}(\{0, (-1, 1)\}) = R_{\mathcal{B}}(\{0, 1/2(-2, 2), 1/3(-3, 3)\})!$$

It is no infimum of purely binary rejection functions.

Weak extension

Setting We have two random variables X and Y , taking values in the finite possibility spaces \mathcal{X} and \mathcal{Y} respectively. From here on, the set of all gambles on $\mathcal{X} \times \mathcal{Y}$ is denoted by \mathcal{L} . This is heavily inspired on [Gert de Cooman & Enrique Miranda, Irrelevant and independent natural extension for sets of desirable gambles].

Gambles: cylindrical extension Let f be a gamble on \mathcal{X} . Define its **cylindrical extension** f^* :

$$f^*(x, y) := f(x) \text{ for all } (x, y) \text{ in } \mathcal{X} \times \mathcal{Y}.$$

f^* belongs to \mathcal{L} . Similarly, for any set A of gambles on \mathcal{X} , we let $A^* := \{f^* : f \in A\}$.

Marginalisation Consider any rejection function R on $\mathcal{X} \times \mathcal{Y}$. Define its **X-marginal** $\text{marg}_X(R)$ as

$$(\text{marg}_X(R))(A) := R(A^*) \text{ for all } A \text{ in } \mathcal{Q}(\mathcal{X}).$$

If R is coherent, then so is $\text{marg}_X(R)$.

Rejection function: weak extension Let R be a coherent rejection function on \mathcal{X} .

What is the least informative coherent rejection function on $\mathcal{X} \times \mathcal{Y}$ that marginalises to R ?

Proposition. The least informative coherent rejection function on $\mathcal{X} \times \mathcal{Y}$ that marginalises to R is $R_{\mathcal{A}}$, where

$$\mathcal{A} := \{A^* : A \in \mathcal{Q}(\mathcal{X}), 0 \in R(A)\}.$$

$R_{\mathcal{A}}$ is called the **weak extension** of R .

Irrelevant natural extension

Conditioning Consider any rejection function R on $\mathcal{X} \times \mathcal{Y}$. For every y in \mathcal{Y} , define its **conditioned rejection function** $R|_y$ on \mathcal{X} as

$$R|_y(A) := \{f \in A : \mathbb{I}_{\{y\}} f \in R(\mathbb{I}_{\{y\}} A)\} \text{ for all } A \text{ in } \mathcal{Q}(\mathcal{X}),$$

where we let $\mathbb{I}_{\{y\}} := \{\mathbb{I}_{\{y\}} f : f \in A\}$ be a set of gambles on $\mathcal{X} \times \mathcal{Y}$. If R is coherent, then so is $R|_y$.

Epistemic irrelevance We say that X is **epistemic irrelevant** to Y when learning the value of X does not influence our beliefs about Y . A rejection function R on $\mathcal{X} \times \mathcal{Y}$ **satisfies epistemic irrelevance** of X to Y when $\text{marg}_Y(R|_x) = \text{marg}_Y(R)$ for all x in \mathcal{X} .

Proposition. Consider any coherent rejection function R on $\mathcal{X} \times \mathcal{Y}$. Then R satisfies epistemic irrelevance of X to Y if and only if

$$(\forall A \in \mathcal{Q}(\mathcal{X})) (\forall y \in \mathcal{Y}) 0 \in R(A) \Leftrightarrow 0 \in R(\mathbb{I}_{\{y\}} A).$$

Let R be a coherent rejection function on \mathcal{X} .

What is the least informative coherent rejection function on $\mathcal{X} \times \mathcal{Y}$ that marginalises to R and satisfies epistemic irrelevance from X to Y ?

Theorem 3. The least informative rejection function on $\mathcal{X} \times \mathcal{Y}$ that marginalises to R and satisfies epistemic irrelevance of X to Y is $R_{\mathcal{A}_{X \rightarrow Y}}$, where

$$\mathcal{A}_{X \rightarrow Y} := \{\mathbb{I}_{\{y\}} A : A \in \mathcal{Q}(\mathcal{X}), 0 \in R(A), y \in \mathcal{Y}\} \cup \{A^* : A \in \mathcal{Q}(\mathcal{X}), 0 \in R(A)\}.$$