

Probability Measures in Gödel $_{\Delta}$ Logic

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- Gödel Logics and Algebras;
 - Gödel $_{\Delta}$ Logic and Algebras;
 - Free Gödel $_{\Delta}$ Algebras;
- States over Free Gödel $_{\Delta}$ Algebras;
- Combinatorial Characterisation of States;
- Adaptation/Generalizations.

Gödel logic G can be semantically defined as a many-valued logic.

Let FORM be the set of formulas over propositional variables x_1, x_2, \dots in the language $\vee, \wedge, \rightarrow, \neg, \perp$.

An assignment is a function

$\mu : \text{FORM} \rightarrow [0, 1] \subseteq \mathbb{R}$ with values in the real unit interval such that, for any two $\alpha, \beta \in \text{FORM}$,

$$\mu(\alpha \wedge \beta) = \min\{\mu(\alpha), \mu(\beta)\},$$

$$\mu(\alpha \vee \beta) = \max\{\mu(\alpha), \mu(\beta)\},$$

$$\mu(\alpha \rightarrow \beta) = \begin{cases} 1 & \text{if } \mu(\alpha) \leq \mu(\beta) \\ \mu(\beta) & \text{otherwise} \end{cases}$$

$$\mu(\neg\alpha) = \mu(\alpha \rightarrow \perp),$$

$$\mu(\perp) = 0,$$

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$$\mu(\Delta(\alpha)) = \begin{cases} 1 & \text{if } \mu(\alpha) = 1 \\ 0 & \text{otherwise} \end{cases}$$

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A **tautology** is a formula α such that $\mu(\alpha) = 1$ for every assignment μ (denoted $\vDash \alpha$).

We write $\vdash \alpha$ to mean that α is derivable from the axioms of G_Δ using *modus ponens* as the only deduction rule.

G_Δ is complete with respect to the many-valued semantics defined above: in symbols, $\vdash \alpha$ if and only if $\vDash \alpha$.

Gödel algebras are Heyting algebras (=Tarski-Lindenbaum algebras of intuitionistic propositional calculus) satisfying the prelinearity equation:

$$(x \rightarrow y) \vee (y \rightarrow x) = \top$$

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An **MTL algebra** $\mathbf{A} = (A, \wedge, \vee, \odot, \rightarrow, \perp, \top)$ is a commutative integral bounded residuated lattice satisfying the **prelinearity** equation,

$$(x \rightarrow y) \vee (y \rightarrow x) = \top$$

A **Gödel Algebra** $\mathbf{A} = (A, \wedge, \vee, \rightarrow, \perp, \top)$ is an **idempotent** MTL Algebra.

The variety \mathbb{G}_Δ is axiomatised as follows,

$$\begin{aligned} \Delta(x) \sqcup \neg\Delta(x) = 1, & \quad \Delta(x \sqcup y) \Rightarrow (\Delta(x) \sqcup \Delta(y)) = 1, & \quad \Delta(x) \Rightarrow x = 1, \\ \Delta(x) \Rightarrow \Delta(\Delta(x)) = 1, & \quad \Delta(x \Rightarrow y) \Rightarrow (\Delta(x) \Rightarrow \Delta(y)) = 1. \end{aligned}$$

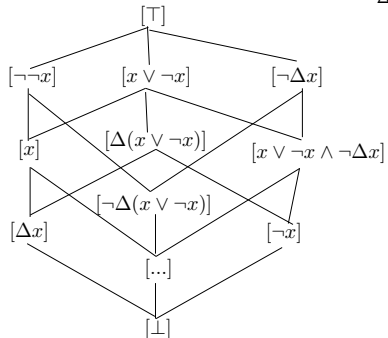
Free Algebras

As usual, $\varphi, \psi \in \text{FORM}_n$ are called **logically equivalent** $\varphi \equiv \psi$, if both $\vdash \varphi \rightarrow \psi$ and $\vdash \psi \rightarrow \varphi$ hold.

The quotient set FORM_n / \equiv endowed with operations $\wedge, \vee, \rightarrow, \Delta, \top, \perp$ induced from the corresponding logical connectives becomes a Gödel $_{\Delta}$ algebra with top and bottom element \top and \perp , respectively.

The specific Gödel $_{\Delta}$ algebra $\mathcal{G}_{\Delta}^n = \text{FORM}_n / \equiv$ is, by construction, the **Lindenbaum algebra** of G_{Δ} over the language $\{x_1, \dots, x_n\}$.

The free 1-generated Gödel $_{\Delta}$ algebra \mathcal{G}_{Δ}^1 :



Lindenbaum algebras are isomorphic to free algebras, and then \mathcal{G}_{Δ}^n is the free n -generated Gödel $_{\Delta}$ algebra \mathbf{F}_n^{Δ} .

Since the variety of Gödel $_{\Delta}$ algebras is locally finite, every finite Gödel $_{\Delta}$ algebra can be obtained as a quotient of a free n -generated Gödel $_{\Delta}$ algebra.

\mathbf{F}_n^Δ is isomorphic to the subalgebra of the algebra of all functions $f: [0, 1]_\Delta^n \rightarrow [0, 1]_\Delta$ generated by the projection $\bar{x}_i: (t_1, \dots, t_n) \mapsto t_i$, for all $i \in \{1, 2, \dots, n\}$.

We write $\bar{\varphi}$ for the elements of \mathbf{F}_n^Δ .

Functional Representation

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The relation \approx on $[0, 1]^n$ is defined as:

$\mathbf{u} = (u_1, \dots, u_n), \mathbf{v} = (v_1, \dots, v_n) \in [0, 1]^n$

$\mathbf{u} \approx \mathbf{v}$ iff there is a permutation σ of $\{1, \dots, n\}$

and a map $\prec: \{0, \dots, n\} \rightarrow \{<, =\}$ such that

$$0 \prec_0 u_{\sigma(1)} \prec_1 \dots \prec_{n-1} u_{\sigma(n)} \prec_n 1$$

iff

$$0 \prec_0 v_{\sigma(1)} \prec_1 \dots \prec_{n-1} v_{\sigma(n)} \prec_n 1$$

\approx is an equivalence relation and $[\mathbf{u}]$ is the equivalence class of \mathbf{u} .

$[0, 1]^n / \approx$ is hence a partition of $[0, 1]^n$.

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With each class $[\mathbf{u}]$, where

$0 \prec_0 u_{\sigma(1)} \prec_1 \dots \prec_{n-1} u_{\sigma(n)} \prec_n 1$,

we associate a unique **ordered partition**

$\rho_{\mathbf{u}} = Q_1 < \dots < Q_h$ of the set

$\{\perp, x_1, \dots, x_n, \top\}$ in the following way:

- $\perp \in Q_1; \top \in Q_h; h > 1;$
- if \prec_i is = then $x_{\sigma(i)}, x_{\sigma(i+1)} \in Q_j;$
- if \prec_i is < and $x_{\sigma(i)} \in Q_j$ then $x_{\sigma(i+1)} \in Q_{j+1}.$

Ordered partitions are in bijections with equivalence classes $[\mathbf{u}] \in [0, 1]^n / \approx$.

When $\rho = \rho_{\mathbf{u}}$, we denote by D_ρ the associated equivalence class $[\mathbf{u}]$.

We write Ω_n for the set of all ordered partitions.

Functional Representation

\mathbf{F}_n^Δ is isomorphic to the subalgebra of the algebra of all functions $f: [0, 1]_\Delta^n \rightarrow [0, 1]_\Delta$ generated by the projection $\bar{x}_i: (t_1, \dots, t_n) \mapsto t_i$, for all $i \in \{1, 2, \dots, n\}$.

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\approx is an equivalence relation and $[\mathbf{u}]$ is the equivalence class of \mathbf{u} .

$[0, 1]^n / \approx$ is hence a partition of $[0, 1]^n$.

A n -variate G_Δ -function is a function $f: [0, 1]^n \rightarrow [0, 1]$ such that for every $\mathbf{u} \in [0, 1]^n$ (equivalently, for any $\rho \in \Omega_n$) the restriction of f to $[\mathbf{u}]$ (equivalently, to D_ρ) is either equal to 0, or to 1, or to a projection function \bar{x}_i .

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we associate a unique **ordered partition**

$$\rho_{\mathbf{u}} = Q_1 < \dots < Q_h$$

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- $\perp \in Q_1; \top \in Q_h; h > 1;$
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Theorem

The elements of \mathbf{F}_n^Δ are exactly the n -variate G_Δ -functions.

Functional Representation

Let $x \triangleleft y = \Delta(x \rightarrow y) \wedge \neg \Delta(y \rightarrow x)$.

Interpreted in $[0, 1]$ we have

$\overline{x \triangleleft y} = 1$ if $x < y$ and $\overline{x \triangleleft y} = 0$ otherwise.

For any $\rho = \rho_{\mathbf{u}} \in \Omega_n$, consider the formula

$$\chi_{\rho} = \bigwedge_{i=0}^n \delta_i,$$

$$\delta_i = \begin{cases} \Delta(x_{\sigma(i)} \leftrightarrow x_{\sigma(i+1)}) & \text{iff } \prec_i \text{ is } =, \\ x_{\sigma(i)} \triangleleft x_{\sigma(i+1)} & \text{iff } \prec_i \text{ is } <. \end{cases}$$

Then it is straightforward to check that

$\overline{\chi_{\rho}(\mathbf{v})} = 1$ iff $\mathbf{v} \approx \mathbf{u}$, while $\overline{\chi_{\rho}(\mathbf{v})} = 0$ otherwise.

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For $n = 2$, the set of Gödel partitions Ω_2 is:

$$\rho_1 = \{0, x, y\} < \{1\}$$

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$$\overline{\chi_{\rho_1}(\rho_1)} = 1$$

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because $\Delta(x \leftrightarrow y) = \Delta(x) = 0$ on ρ_5 .

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Let $f: [0, 1]^n \rightarrow [0, 1]$ be a G_{Δ} -function.

y_{ρ} is the element of $\{\perp, x_1, \dots, x_n, \top\}$ such that $\overline{y_{\rho}}$ coincides with f over the whole of D_{ρ} .

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$$\varphi = \bigvee_{\rho \in \Omega_n} (\chi_{\rho} \wedge y_{\rho}).$$

For any point $\mathbf{u} \in [0, 1]^n$, $\overline{\varphi(\mathbf{u})}$ coincides with $\overline{\chi_{\rho} \wedge y_{\rho}(\mathbf{u})}$ for the unique $\rho \in \Omega_n$ such that $\mathbf{u} \in D_{\rho}$.

Then, $\overline{\varphi} = f$.

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A **state** on \mathbf{F}_n^Δ is a function $s: \mathbf{F}_n^\Delta \rightarrow [0, 1]$ such that, for every $f, g \in \mathbf{F}_n^\Delta$:

- 1 $s(\perp) = 0, s(\top) = 1$;
- 2 $s(f \vee g) = s(f) + s(g) - s(f \wedge g)$;
- 3 If $f \leq g$ then $s(f) \leq s(g)$;
- 4 If $f \leq g$ and $s(g) = s(f)$ then $s(\Delta(g \rightarrow f)) = 1$.

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- ④ If $f \leq g$ and $s(g) = s(f)$ then $s(\Delta(g \rightarrow f)) = 1$.

Theorem

The following hold.

- ① If $s: \mathbf{F}_n^\Delta \rightarrow [0, 1]^n$ is a state, there exists a Borel probability measure μ on $[0, 1]^n$ such that

$$\int_{[0,1]^n} f \, d\mu = s(f), \text{ for every } f \in \mathbf{F}_n^\Delta.$$

- ② Viceversa, for any Borel probability measure μ on $[0, 1]^n$, the function $s: \mathbf{F}_n^\Delta \rightarrow [0, 1]$ defined by the above integral is a state.

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States of \mathbf{F}_n^{Δ} are the convex combinations of finitely many truth value assignments.

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Corollary

States of \mathbf{F}_n^{Δ} are the convex combinations of finitely many truth value assignments.

Let s be the state on $\mathbf{F}_2(\mathbb{G}_{\Delta})$ given by

$$s(\overline{\chi_{\rho_1}}) = 1/3 \qquad s(\overline{\chi_{\rho_4}}) = 1/6$$

$$s(\overline{\chi_{\rho_5}}) = 1/2 \qquad s(\overline{x \wedge \chi_{\rho_4}}) = 1/12$$

$$s(\overline{x \wedge \chi_{\rho_5}}) = 1/12 \qquad s(\overline{y \wedge \chi_{\rho_4}}) = 1/6$$

$$s(\overline{\chi_{\sigma}}) = 0 \qquad \text{for } \sigma \notin \{\rho_1, \rho_4, \rho_5\}$$

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Define the discrete measure μ by setting

$$\mu(\{\mathbf{z}_{\rho_1}\}) = 1/3 \qquad \mathbf{z}_{\rho_1} = (0, 0)$$

$$\mu(\{\mathbf{z}_{\rho_4}\}) = 1/6 \qquad \mathbf{z}_{\rho_4} = (1/2, 1/2)$$

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Corollary

States of \mathbf{F}_n^{Δ} are the convex combinations of finitely many truth value assignments.

Let s be the state on $\mathbf{F}_2(\mathbb{G}_{\Delta})$ given by

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$$\begin{aligned} s(f) &= s(\overline{\chi_{\rho_1}} \vee (\bar{y} \wedge \overline{\chi_{\rho_5}})) = s(\overline{\chi_{\rho_1}}) + s(\bar{y} \wedge \overline{\chi_{\rho_5}}) = \frac{1}{3} + \frac{1}{6} = \frac{1}{2} \\ \int_{[0,1]^2} f d\mu &= \sum_{i \in \{1,4,5\}} f(z_{\rho_i}) \mu(\{z_{\rho_i}\}) = 1 \cdot \frac{1}{3} + 0 \cdot \frac{1}{6} + \frac{1}{3} \cdot \frac{1}{2} = \frac{1}{2} = s(f). \end{aligned}$$

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A nonempty subset F of A is called an *upper-set* when for all $x, y \in A$, if $x \leq y$ and $x \in F$, then $y \in F$. If $x \odot y \in F$ for all $x, y \in F$, then F is a **filter** of \mathbf{A} . We call $\bigwedge_{x \in F} x$ the *generator* of the filter F . A filter F of A is **prime** if $F \neq A$ and for all $x, y \in A$, either $x \rightarrow y \in F$ or $y \rightarrow x \in F$.

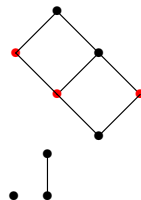
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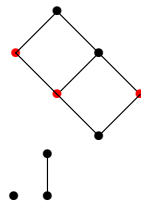
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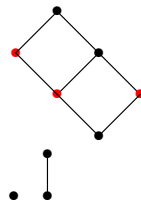
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Proposition (Aguzzoli and Codara, 2016)

Every finite G_Δ -algebra A is a direct product of chains. That is, $A \simeq \prod_{F \in \text{Max}(A)} A/F$, and $\text{Max}(A) = \text{Spec}(A)$.

For each $\mathbf{A} \in (\mathbb{G}_{\Delta})_{fin}$, the poset $Spec(\bar{\mathbf{A}})$, that is, the prime spectrum of the G -algebra reduct of \mathbf{A} , ordered by reverse inclusion, is isomorphic with the poset of the j.i. elements of \mathbf{A} .

$$Spec^{\Delta}(\mathbf{A}) = \mathcal{C}(Spec(\bar{\mathbf{A}}))$$

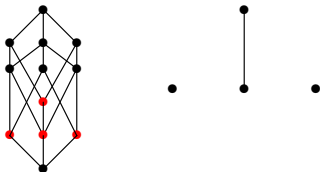
where $\mathcal{C}(P)$ is the multiset $\{C_1, C_2, \dots, C_u\}$, when the poset P is a disjoint union $C_1 \cup C_2 \cup \dots \cup C_u$ of chains.

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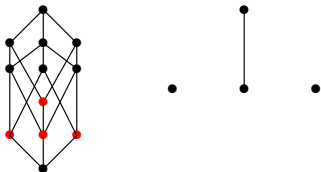


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Conversely, given a chain C we define:

$$Sub^{\Delta}(C) = \{D \subseteq C \mid D = \downarrow D\}, \cup, \cap, \rightarrow, \sim, \emptyset, C, \Delta,$$

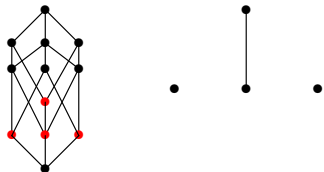
where $\Delta C = C$ and $\Delta D = \emptyset$, for each subchain $D \subsetneq C$,
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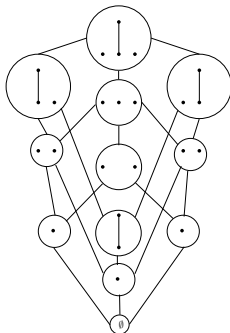
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MC be the category whose objects are finite multisets of (nonempty) finite chains, and whose morphisms $h: C \rightarrow D$, are defined as follows.

Display C as $\{C_1, \dots, C_m\}$ and D as $\{D_1, \dots, D_n\}$. Then $h = \{h_i\}_{i=1}^m$, where each h_i is an order preserving surjection $h_i: C_i \twoheadrightarrow D_j$ for some $j \in \{1, 2, \dots, n\}$.

$(\mathbb{G}_\Delta)_{fin}$ is the full subcategory of \mathbb{G}_Δ whose objects have finite cardinality, and morphisms are simply homomorphisms of algebras.

The two previous constructions are functorial:

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Theorem (Aguzzoli and Codara, 2016)

The categories $(\mathbb{G}_{\Delta})_{fin}$ and MC are dually equivalent.

A **labeling** l is a function $l: \text{Spec}^\Delta \mathbf{F}_n^\Delta \rightarrow [0, 1]$, such that

- 1 $\sum_{p \in \text{Spec}^\Delta \mathbf{F}_n^\Delta} l(p) = 1$;
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Theorem

Let S_n be the collection of all states $s: \mathbf{F}_n^\Delta \rightarrow [0, 1]$, and let L_n be the collection of all labelings $l: \text{Spec}^\Delta \mathbf{F}_n^\Delta \rightarrow [0, 1]$. Then, the map defined for every formula φ over the set of variables $\{x_1, \dots, x_n\}$ by

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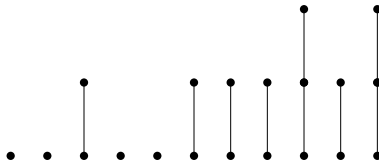
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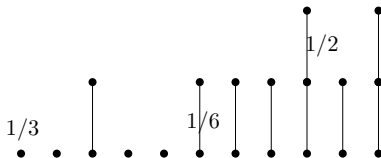
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Further axiomatisations

Drastic Product algebras constitute the subvariety \mathbb{DP} of MTL axiomatised by $x \sqcup \sim(x * x) = 1$.

Theorem (Aguzzoli, Bianchi and V., 2014)

\mathbb{MC}^\top is dually equivalent to the category \mathbb{DP}_{fin} of finite DP algebras and their homomorphisms.

\mathbb{MC}^\top is a non-full subcategory of MC.

Hence, we can adapt the presented constructions to axiomatise States over DP algebras.

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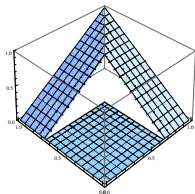
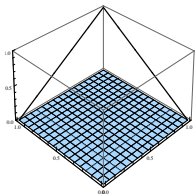
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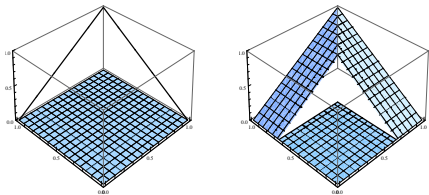
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